

# A new method for calculating the chromatic polynomial

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**Abstract.** Chromatic polynomials are widely used in graph theoretical or chemical applications in many areas. Birkhoff-Lewis theorem is the most important tool to find the chromatic polynomial of any given graph. Here we obtain several shortcut moves to calculate this polynomial covering all graphs.

**M.S.C. 2010:** 05C15, 05C31, 05C10, 05C30.

**Key words:** chromatic polynomial; chromatic number; graph colouring.

## 1 Introduction

One of the well-known applications of graph theory is the 4-colour problem. There are many notions related to colourings of graphs. A colouring of a graph  $G(V, E)$  is a mapping  $f : V \rightarrow C$ , where  $C$  is the set of colours, with  $f(u) \neq f(v)$  for  $uv \in E$ .

If there is a colouring of  $G$  with  $n$  colours, then  $G$  is said to be  $n$ -colourable. The smallest number  $n$  for which  $G$  is  $n$ -colourable is called the chromatic number and denoted by  $\chi(G)$ . Clearly a graph with at least one loop cannot have any colourings. As we shall be considering simple graphs, this case will be out of question. Similarly, edge colourings are defined.

Here we shall study another aspect related to colourings, the chromatic polynomial of a graph. The chromatic polynomial of  $G$  is defined to be a function  $C_G(k)$  which expresses the number of distinct  $k$ -colourings possible for the graph  $G$  for each integer  $k > 0$ . This number was first used by Birkhoff in 1912. Chromatic polynomials are widely used in determining several properties of graphs. Although there is no known formula for the chromatic polynomial of any given graph, there are algorithms to do that. The most well-known algorithm is the Birkhoff-Lewis theorem stated as below:

**Theorem 1.1** (Birkhoff-Lewis, 1946). *The chromatic polynomial of a graph  $G$  can be found by the formula*

$$C_G(k) = C_{G-\epsilon}(k) - C_{G/\epsilon}(k),$$

where  $G - e$  is the graph obtained by deleting the edge  $e$  from  $G$ , and  $G/e$  is the graph obtained from  $G$  by removing  $e$  and identifying the end vertices of  $e$  and leaving only one copy of any resulting multiple edges (see Fig. 1.1 and 1.2).

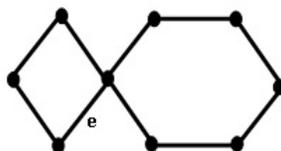


Figure 1.1 A graph  $G$



Figure 1.2 The graphs  $G - e$  and  $G/e$

As usual, we denote path, cycle, star, complete, complete bipartite and tadpole graphs by  $P_n$ ,  $C_n$ ,  $S_n$ ,  $K_n$ ,  $K_{r,s}$  and  $T_{r,s}$ , respectively.

A bridge of a connected graph is an edge whose removal disconnects the graph. In another words, a bridge is an edge of a graph  $G$  whose removal increases the number of components of  $G$ , [5], p. 26. Also it is well known that an edge of a connected graph is a bridge iff it does not lie on any cycle. For the details of these and related notions, see [1], [7], [3], [4] and [6].

## 2 Calculating the chromatic polynomial by splitting the given graph

In this section, we want to obtain a new and alternative method to find the chromatic polynomial of any given graph. The idea behind this will be splitting the given graph to smaller subgraphs. We shall state results for the chromatic polynomial of a given graph  $G$  in terms of two smaller graphs  $G_1$  and  $G_2$ . If necessary, these can be generalized to a finite number of subgraphs. We shall consider four different cases where a graph is splitted into two parts at one vertex or through an edge which is called a bridge.

In this paper we use  $G_1 \cup G_2$  to denote the union of two graphs  $G_1$  and  $G_2$  which have one joint vertex.

One can think of starting by any graph and applying Theorem 1.1 to get to a point eventually. But these two theorems cannot be always enough for all graphs. Some more complex graphs need other splitting methods.

We first have the following useful result:

**Theorem 2.1.** *Let  $G_1$  and  $G_2$  be two graphs with no common vertex. Let also  $G_1 \coprod G_2$  be the disjoint union of  $G_1$  and  $G_2$ . Then*

$$C_G(k) = C_{G_1}(k) \cdot C_{G_2}(k).$$

*Proof.* As  $G_1$  and  $G_2$  are disjoint graphs, the colouring of  $G_1$  is independent from the colouring of  $G_2$ . By the principal of counting, the result follows.  $\square$

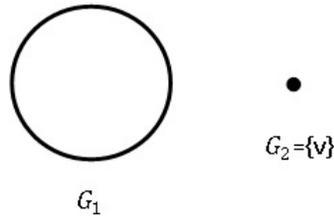
This result can be generalized to the product of a finite number of graphs by mathematical induction:

**Corollary 2.2.** *Let  $G_1, G_2, \dots, G_n$  be pairwise disjoint graphs and also let  $G = \coprod_{i=1}^n G_i$  be their disjoint union. Then*

$$C_G(k) = \prod_{i=1}^n C_{G_i}(k).$$

We now give a result which will be needed later:

**Corollary 2.3.** *Let  $G$  be a graph with two components, one being a unique vertex, see Figure 1.3.*



**Figure 1.3** A graph  $G = G_1 \coprod \{v\}$

Then

$$C_G(k) = k \cdot C_{G_1}(k).$$

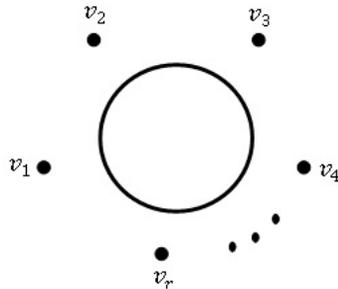
*Proof.* As the colouring of  $G_2 = \{v\}$  is independent from the colouring of  $G_1$ , the result follows.  $\square$

This easily can be generalized as follows:

**Corollary 2.4.** *Let  $G$  have  $r$  isolated vertices. That is*

$$G = G_1 \coprod \{v_1\} \coprod \cdots \coprod \{v_r\},$$

see Figure 1.4.

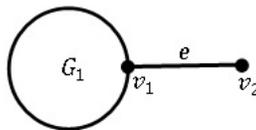


**Figure 1.4** A graph  $G = G_1 \coprod \{v_1\} \coprod \cdots \coprod \{v_r\}$  with  $r$  isolated vertices

Then

$$C_G(k) = k^r \cdot C_{G_1}(k).$$

Next, we consider the case where  $G$  has only one pendant vertex  $v$ . Let us term the edge connecting  $v$  to the rest of the graph by  $e$ , and let  $u$  be the vertex of  $G$  connected to  $v$  by  $e$ , see Figure 1.5.



**Figure 1.5** The graph  $G$  with one pendant vertex

Then we have the following relation between the chromatic polynomials of  $G$  and  $G_1$ :

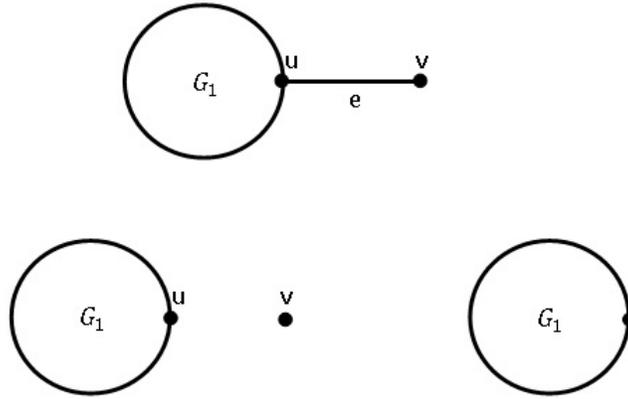
**Lemma 2.5.** *Let  $G$  have only one pendant vertex  $v$ . Then*

$$C_G(k) = (k - 1) \cdot C_{G_1}(k).$$

*Proof.* We shall use the Birkhoff-Lewis Theorem. Recall that

$$C_G(k) = C_{G-e}(k) - C_{G/e}(k),$$

see Figure 1.6.



**Figure 1.6** The graph  $G$ ,  $G - e$  and  $G/e$

By Lemma 2.1, we can write

$$C_{G-e}(k) = k \cdot C_{G_1}(k).$$

Further

$$C_{G/e}(k) = C_{G_1}(k)$$

implies that

$$\begin{aligned} C_G(k) &= k \cdot C_{G_1}(k) - C_{G_1}(k) \\ &= (k - 1) \cdot C_{G_1}(k). \end{aligned}$$

□

Let  $G$  be a graph having  $r$  pendant vertices. See Figure 1.7.

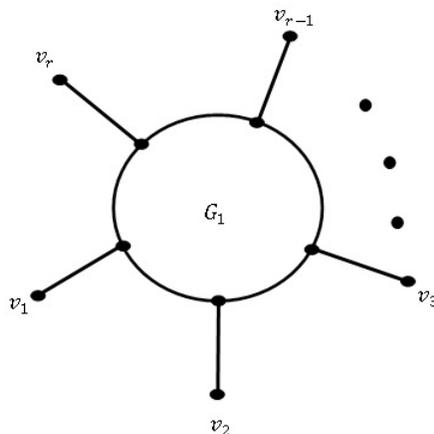


Figure 1.7 The graph  $G$  with  $r$  pendant vertices

Many graphs have pendant vertices and it is useful to be able to calculate the chromatic polynomial of the remaining graph  $G_1$  in Fig. 1.7 after omitting all pendant vertices and edges connecting them to  $G_1$ . To achieve this, we have the following generalisation for a graph  $G$  with  $r$  pendant vertices:

**Lemma 2.6.** *Let  $G$  have  $r$  pendant vertices. Then*

$$C_G(k) = (k - 1)^r \cdot C_{G_1}(k).$$

*Proof.* It follows from the proof of Lemma 2.5. □

**Theorem 2.7.** *Let  $G_1$  be a graph and  $G$  be the graph obtained by joining  $G_1$  with a path graph  $P_r$  as in Figure 1.8.*

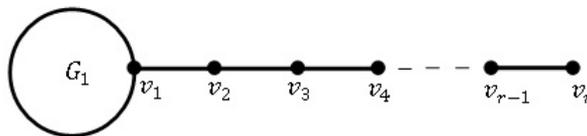
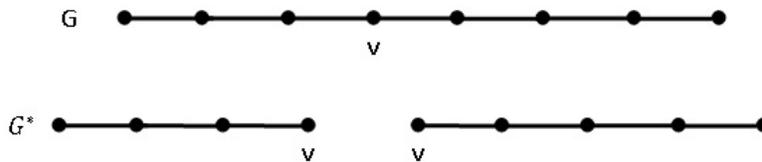


Figure 1.8 The graph  $G = G_1 \cup P_r$

Then

$$C_G(k) = \frac{C_{G_1}(k) \cdot C_{P_r}(k)}{k}.$$

**Example 2.1.** If we choose the graphs  $P_4$  and  $P_5$  joined at a vertex  $v$ ,



**Figure 1.9** The graph  $G = P_8$ ,  $G^* = P_4 \amalg P_4$

We know that

$$C_G(k) = C_{P_8}(k) = k \cdot (k-1)^7.$$

Also by Theorem 2.1, we have

$$C_{G^*}(k) = C_{P_4}(k) \cdot C_{P_4}(k).$$

Therefore

$$\begin{aligned} C_{G^*}(k) &= k \cdot (k-1)^3 \cdot k \cdot (k-1)^4 \\ &= k^2 \cdot (k-1)^7. \end{aligned}$$

Then

$$\frac{C_G(k)}{C_{G^*}(k)} = \frac{k \cdot (k-1)^7}{k^2 \cdot (k-1)^7} = \frac{1}{k}$$

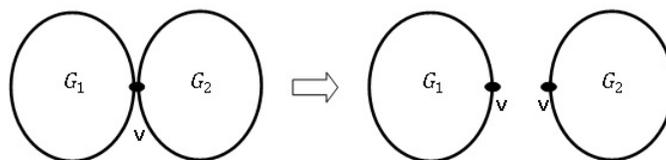
so

$$C_{G^*}(k) = \frac{C_{P_8}(k) \cdot C_{P_4}(k)}{k}.$$

Now we deal with the graphs which are connected at a cut vertex:

**Lemma 2.8.** *Let the graphs  $G_1$  and  $G_2$  be connected at a vertex  $v$  as in Figure 1.10. Then*

$$C_G(k) = \frac{C_{G_1}(k) \cdot C_{G_2}(k)}{k}.$$

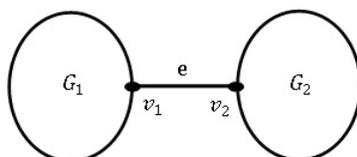


**Figure 1.10** First way of splitting a graph  $G = G_1 \cup G_2$

**Theorem 2.9.** Let the graph  $G$  be the union of graphs  $G_1$  and  $G_2$  connected by a bridge  $e$ , as in Figure 1.11. Then

$$C_G(k) = \frac{C_{G_1}(k) \cdot C_{G_2}(k)}{k \cdot (k - 1)}.$$

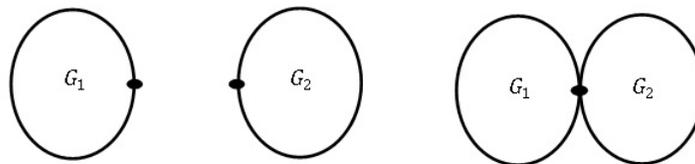
*Proof.* Let the graph  $G$  be the union of graphs  $G_1$  and  $G_2$  connected by a bridge  $e$ . See Figure 1.11.



**Figure 1.11** The graphs  $G_1$  and  $G_2$  connected by a bridge  $e$

By Birkhoff-Lewis theorem we have

$$C_G(k) = C_{G-e}(k) - C_{G/e}(k).$$



**Figure 1.12** The graphs  $G - e$  and  $G/e$

By Theorem 2.1,

$$C_{G-\epsilon}(k) = C_{G_1}(k) \cdot C_{G_2}(k)$$

and by Lemma 2.3, we have

$$k \cdot C_{G/\epsilon}(k) = C_{G_1}(k) \cdot C_{G_2}(k).$$

Therefore we obtain

$$C_{G/\epsilon}(k) = \frac{C_{G_1}(k) \cdot C_{G_2}(k)}{k}.$$

If we substitute these in

$$C_G(k) = C_{G-\epsilon}(k) - C_{G/\epsilon}(k),$$

we obtain

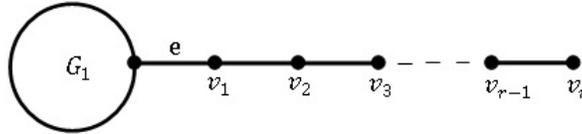
$$\begin{aligned} C_G(k) &= C_{G_1}(k) \cdot C_{G_2}(k) - \frac{C_{G_1}(k) \cdot C_{G_2}(k)}{k} \\ &= C_{G_1}(k) \cdot C_{G_2}(k) \cdot \left(1 - \frac{1}{k}\right) \\ &= \frac{C_{G_1}(k) \cdot C_{G_2}(k) \cdot (k-1)}{k} \end{aligned}$$

which is the required result.  $\square$

The following result gives an important case:

**Corollary 2.10.** *Let a graph  $G$  be the union of a graph  $G_1$  and a path  $P_r$  connected by a bridge  $e$ . Then*

$$C_G(k) = C_{G_1}(k) \cdot (k-1)^r.$$



**Figure 1.13** A graph  $G_1$  and a path  $P_r$  connected by a bridge  $e$

*Proof.* By Lemma 2.3, we have

$$C_G(k) = \frac{C_{G_1}(k) \cdot C_{P_r}(k) \cdot (k-1)}{k}.$$

Also as

$$C_{P_r}(k) = k \cdot (k - 1)^{r-1},$$

we obtain

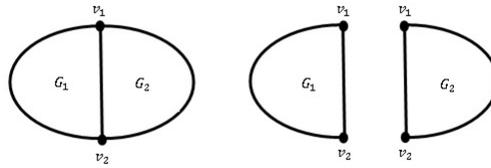
$$C_G(k) = C_{G_1}(k) \cdot (k - 1)^r.$$

□

Now we deal with another case where we split a graph into two subgraphs through a common edge as in Figure 1.14.

**Lemma 2.11.** *Let  $G$  be a graph which can be splitted into two subgraphs  $G_1$  and  $G_2$  through a common edge  $e = v_1v_2$ . Then*

$$C_G(k) = \frac{C_{G_1}(k) \cdot C_{G_2}(k)}{k \cdot (k - 1)}.$$



**Figure 1.14** A graph  $G$  splitted through an edge

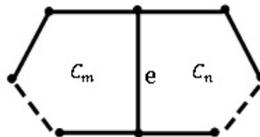
*Proof.* It follows by Birkhoff Lewis theorem. □

Now we deal with another common case where a graph  $G$  can be splitted into two cycle graphs  $C_m$  and  $C_n$  through a common edge  $e$ :

**Theorem 2.12.** *Let  $G$  be a graph which can be splitted into two cycle graphs  $C_m$  and  $C_n$  through a common edge  $e$  as in Figure 1.15. Then we have*

$$C_G(k) = \frac{C_{C_m}(k) \cdot C_{C_n}(k)}{k \cdot (k - 1)}.$$

*Proof.* The given graph  $G$  is as in the following figure:



**Figure 1.15** Splitting  $G$  into two cycle graphs  $C_m$  and  $C_n$  through a common edge  $e$

By Birkhoff-Lewis theorem,



**Figure 1.16** The graphs  $G - e$  and  $G/e$

$$G - e = C_{m+n-2}$$

and by Lemma 2.3, we have

$$C_{G/e}(k) = \frac{C_{C_{m-1}}(k) \cdot C_{C_{n-1}}(k)}{k}$$

$$\begin{aligned} C_G(k) &= \frac{(k-1)^{m+n-2} + (-1)^{m+n-2} \cdot (k-1)}{[(k-1)^{m-1} + (-1)^{m-1} \cdot (k-1)] \cdot [(k-1)^{n-1} + (-1)^{n-1} \cdot (k-1)]} \\ &= \frac{k \cdot (k-1)^{m+n-2} + k \cdot (-1)^{m+n-2} \cdot (k-1) - (k-1)^{m+n-2}}{k} \\ &= \frac{(-1)^{n-1} \cdot (k-1)^m + (-1)^{m-1} \cdot (k-1)^n + (-1)^{m+n-2} \cdot (k-1)^2}{k} \\ &= \frac{(k-1)^{m+n-2} \cdot (k-1) + (-1)^{m+n-2} \cdot (k-1) + (-1)^{n-1} \cdot (k-1)^m}{k} \\ &= \frac{(-1)^{m-1} \cdot (k-1)^n}{k} \\ &= \frac{(k-1)^{m+n-1} + (-1)^{m+n-2} \cdot (k-1) + (-1)^{n-1} \cdot (k-1)^m}{k} \\ &= \frac{(-1)^{m-1} \cdot (k-1)^n}{k} \end{aligned}$$

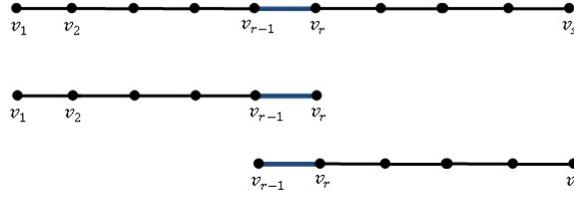
and as a consequence, we get the required result:

$$C_{G(k)} = \frac{C_{C_m(k)} \cdot C_{C_n(k)}}{k \cdot (k-1)}. \quad \square$$

**Theorem 2.13.** Let  $1 < r < s$ . For a graph  $G$  which is obtained by overlapping one edge of two path graphs  $P_r$  ve  $P_{s-r+2}$ , we have

$$C_G(k) = \frac{C_{P_r}(k) \cdot C_{P_{s-r+2}}(k)}{k \cdot (k-1)}.$$

*Proof.* Let  $G$  be the graph in the first line of Figure 1.17. Let us separate  $G$  into two paths  $P_r$  ve  $P_{s-r+2}$  at the edge  $v_{r-1}v_r$  as in the second line of Figure 1.17.



**Figure 1.17** The graphs  $G, P_r, P_{s-r+2}$

From the formula of chromatic polynomial of the path graph, we have

$$C_G(k) = C_{P_s}(k) = k \cdot (k - 1)^{s-1}, \quad C_{P_r}(k) = k \cdot (k - 1)^{r-1},$$

$$C_{P_{s-r+2}}(k) = k \cdot (k - 1)^{s-r+1},$$

and hence we conclude that

$$C_{P_r}(k) \cdot C_{P_{s-r+2}}(k) = k \cdot (k - 1)^{r-1} \cdot k \cdot (k - 1)^{s-r+1} = k^2 \cdot (k - 1)^s.$$

This gives us the required result:

$$C_G(k) = \frac{C_{P_r}(k) \cdot C_{P_{s-r+2}}(k)}{k \cdot (k - 1)}. \quad \square$$

**Acknowledgements.** The second author is supported by the Research Fund of Uludag University, Project No: 2015-17.

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